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## Datta et al.

#### (54) ADAPTIVE THRESHOLDS FOR HIGH SPEED DOWNLINK SHARED CONTROL CHANNEL (HS-SCCH) (PART I) DETECTION SCHEMES

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- (51) Int. Cl.

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See application file for complete search history.

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### (57) **ABSTRACT**

A technique is provided for implementing adaptive thresholds associated with HS-SCCH detection schemes; and when applied to any HS-SCCH detection scheme, the resulting false alarm probability curves are more robust to amplitude variations of the different shared control channels. The technique has low computational complexity and low storage requirements for the estimator. Such lower complexity detection schemes have been found to outperform more complex schemes when adaptive thresholds are applied.

#### 9 Claims, 2 Drawing Sheets







FIG. 5



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### ADAPTIVE THRESHOLDS FOR HIGH SPEED DOWNLINK SHARED CONTROL CHANNEL (HS-SCCH) (PART I) DETECTION SCHEMES

#### CLAIM TO PRIORITY OF PROVISIONAL APPLICATION

The application claims priority under 35 U.S.C. §119(e)(1) of provisional application Ser. No. 60/392,769, filed Jul. 1, 2002, by Suparna Datta-Bellamy, Anand G. Dabak and Timo- 10 thy M. Schmidl.

#### BACKGROUND OF THE INVENTION

1. Field of the Invention

This invention relates generally to the shared control channel (HS-SCCH) for the HS-DSCH (high speed downlink shared channel) defined in technical standard 3GPP TS 25.212 for multiplexing and channel coding (FDD), developed within the  $3^{rd}$  Generation Partnership Project, and more 20 particularly to a technique for determining adaptive thresholds associated with HS-SCCH detection schemes.

2. Description of the Prior Art

Using fixed threshold values in HS-SCCH detection schemes has been shown to yield false alarm probabilities that 25 increase with signal-to-noise ratio (SNR). Specifically, as the SNR increases, the Viterbi decoder metric increases. Consequently, the received signal is more likely to cross the threshold.

In view of the above, a technique for implementing adap-30 tive thresholds associated with HS-SCCH detection schemes and that is robust to amplitude variations of the different shared control channels would be both advantageous and desirable.

#### SUMMARY OF THE INVENTION

The present invention is directed to a technique for implementing adaptive thresholds associated with HS-SCCH detection schemes; and when applied to any detection 40 scheme, the resulting false alarm probability curves are more robust to amplitude variations of the different shared control channels. The technique has low computational complexity and low storage requirements for the estimator. Such lower complexity detection schemes have been found to outperform 45 more complex schemes when adaptive thresholds are applied.

According to one embodiment, a high-speed shared control channel (HS-SCCH) detection method comprises the steps of estimating an amplitude value associated with a received signal on a shared control channel (SCCH); and 50 scaling a desired threshold value by the estimated amplitude value to generate an adaptive threshold value there from, such that the adaptive threshold value is substantially robust to amplitude variations associated with at least one SCCH.

#### BRIEF DESCRIPTION OF THE DRAWINGS

Other aspects and features of the present invention and many of the attendant advantages of the present invention will be readily appreciated as the invention becomes better understood by reference to the following detailed description when considered in connection with the accompanying drawings in which like reference numerals designate like parts throughout the figures thereof and wherein:

FIG. 1 is a flow diagram illustrating part I coding including 65 a UE ID scrambler associated with the currently approved structure for the HS-SCCH;

FIG. **2** is a timing diagram illustrating the relationship between HS-SCCH and HS-DSCH timing;

FIG. **3** is a block diagram illustrating an adaptive threshold calculation technique associated with four shared control channels according to one embodiment of the present invention:

FIG. **4** is a graph illustrating simulated miss-probability plots associated with a plurality of detection schemes using adaptive threshold values; and

FIG. **5** is a graph illustrating simulated false alarm probability plots associated with a plurality of detection schemes using adaptive threshold values.

While the above-identified drawing figures set forth alternative embodiments, other embodiments of the present invention are also contemplated, as noted in the discussion. In all cases, this disclosure presents illustrated embodiments of the present invention by way of representation and not limitation. Numerous other modifications and embodiments can be devised by those skilled in the art which fall within the scope and spirit of the principles of this invention.

#### DETAILED DESCRIPTION OF THE PREFERRED EMBODIMENTS

The present embodiments are best understood with reference to FIG. 1, by first providing a background for the currently approved structure for the HS-SCCH that is broken up into 2 parts. Part I consists of 8 information bits including the CCS and the modulation level; while Part II contains 13 information bits and a 16-bit CRC. Looking now at FIG. 1, coding 100 for Part I consists of inputting the 8 information bits plus 8 tail bits into a rate  $\frac{1}{3}$  convolutional coder and rate-matching the output to 40 bits as shown in block 102. These encoded output bits are then scrambled by a 16-bit UE

<sup>35</sup> ID, which is extended to 40 bits via a punctured, rate <sup>1/2</sup> convolutional code as shown in blocks **104** and **106**. These scrambled bits are carried entirely in the first slot **202** of the HS-SCCH transmission **200** shown in FIG. **2**.

In Part II (enumerated **204** in FIG. **2**), the CRC is scrambled by the 16-bit UE ID, after which all 29 bits plus 8 tail bits are convolutionally encoded by the rate <sup>1</sup>/<sub>3</sub> code. The resulting 111 output bits are rate-matched to 80 bits and transmitted over 2 slots. Rate-matching for both Part I and Part II is implemented using the following optimal puncturing patterns. For Part I, puncture bits **1**, **2**, **4**, **8**, **42**, **45**, **47**, **48** (where the bit count ranges from 1 to 48). For Part II, puncture bits **1**, **2**, **3**, **4**, **5**, **6**, **7**, **8**, **12**, **14**, **15**, **24**, **42**, **48**, **54**, **57**, **60**, **66**, **69**, **96**, **99**, **101**, **102**, **104**, **105**, **106**, **107**, **108**, **109**, **110**, **111**.

FIG. 2 is a timing diagram illustrating the relationship between HS-SCCH 200 and HS-DSCH 206 timing. The total number of HS-SCCH's 200 can range from one to four, with all the UE's trying to decode Part I of all the SCCH's 200. The decoding process (for each SCCH) consists of first descrambling the bits with its unique ID, and then sending those bits through the convolutional decoder. Because there is no CRC for Part I, each UE needs to have some kind of detection mechanism to determine whether or not to buffer and decode Part II. The performance metrics that this detection mechanism needs to address are: probability of miss ( $p_{miss}$ ) and probability of false alarm ( $p_{fa}$ ).

A miss event occurs when an SCCH transmission **200** is intended for a UE, and yet the UE fails to realize this and does not proceed to decode Part II. This is a serious occurrence that will result in a loss of throughput and wasted resources. Consequently, these types of error events should most preferably be minimized to the order of 1-2%. 15

A false alarm event occurs when a UE incorrectly determines that a transmission was intended for itself and proceeds to decode Part II **204** on that SCCH **200**. The consequence in this case is false buffering and increased power consumption. Note that because Part II has a strong 16-bit UE specific CRC, 5 it is unlikely that this type of error will result in any HARQ (hybrid automatic repeat request) combining loss.

In view of the foregoing, the present inventors then examined six detection schemes, including:

1. Absolute Viterbi Metric (AVM);

2. AVM+1 Re-Encode;

3. Yamamoto-Itoh Algorithm (YI)+Minimum Path Metric Difference (MPMD);

4. AVM+YI;

5. AVM+YI+1 Re-Encode; and

6. Re-Encoded Symbol Error Rate (SER).

The Absolute Viterbi Metric (AVM) has the UE choose the 20 SCCH with the largest Viterbi metric. If that metric is greater than some threshold value, then the UE determines that the SCCH is right for it.

AVM+1 Re-Encode entails first choosing the SCCH with the largest Viterbi metric. Next, re-encode the decoded bits of 25 that SCCH. If less than some threshold number of bits disagrees with the original (hard) received symbols, then the UE determines that the SCCH is right for it.

The basic idea behind the Yamamoto-Itoh algorithm is that whenever two merging paths in the Viterbi decoding trellis are 30 very close in terms of path metric, the survivor path is more likely to be in error. So, the YI algorithm sets a threshold on the path metric difference (PMD) between the two largest merging paths at each state in the trellis. If the difference is less than the threshold, then the surviving path is labeled 35 "bad". If, at the end of Viterbi decoding, any leg of the final surviving path is "bad", then the transmission is declared a failure. In the case where more than one of the four transmissions has a "good" surviving path, the secondary soft metric Minimum Path Metric Difference (MPMD) is used. While 40 decoding, we keep track of the minimum path metric difference for all four channels. If the YI algorithm yields a tie between more than one SCCH, then we choose the channel with the maximum MPMD.

AVM+YI involves first choosing the SCCH with the largest 45 Viterbi metric. If the YI algorithm labels that surviving path "good", then the UE determines that the SCCH is right for it; otherwise the transmission is declared a failure.

AVM+YI+1 Re-Encode requires first completing AVM+ YI. If the SCCH passes the YI test, then perform a re-encode 50 of the decoded bits. If less than some threshold number of bits disagrees with the original (hard) received symbols, then the UE determines that the SCCH is right for it.

The SER metric entails performing the Re-Encode metric for all for four SCCH's. The UE compares the number of 55 errors to a threshold to determine which (if any) of the transmissions was right for it.

While studying the foregoing detection schemes, an issue of concern was discovered by the present inventors to lie with using fixed threshold values. Specifically, the Viterbi metric 60 was found to increase as the SNR increases. The received signal is therefore more likely to cross the threshold, thus leading to a worse false alarm probability at high SNR values. To counter this, the present inventors decided to investigate the use of adaptive thresholds that are based on estimating the 65 (QPSK) amplitude of the received signals. One adaptive technique is now described below as follows:

4

Let the received bits on a shared control channel after despreading be indicated by the vector

$$r = \{r_l^1 + j^* r_Q^1, \dots, r_l^{40} + j^* r_Q^{40}\}$$
 (complex *I*, *Q*)

Noticeably, there are up to four shared control channel vectors similar to the one defined above. Then the received signal <sup>10</sup> is normalized by the estimate

$$\hat{A} = \sqrt{2} * \left( \frac{\sum_{i=1}^{40} |r_i^i| + |r_0^i|}{(2*40)} \right).$$

That is, we generate a new vector  $\tilde{r}=r/\hat{A}$  and pass it to the Viterbi decoder. In this way, the thresholds are now effectively adaptive and are robust to amplitude variations of the different shared control channels. Other normalization algorithms like using the Euclidean metric instead of the absolute can also be used to normalize the amplitude of the received signal. In this case, we would have

$$\hat{A} = \left(\frac{\sqrt{\left(\sum_{i=1}^{40} |r_i^i|^2 + |r_Q^i|^2\right)^2 / 40}}{2}\right)$$

Alternatively, rather than buffering and normalizing the received signal, we can instead compute the estimate  $\hat{A}$  as the received bits enter the Viterbi decoder and then scale the threshold value t by the final result such as seen in FIG. **3** that is a block diagram illustrating an adaptive threshold calculation technique **300** associated with four shared control channels **302 304 306**, **308** according to one embodiment of the present invention. That is, we generate the new adaptive threshold value  $\tilde{t} = t^* \hat{A}$ .

The present inventors ran simulations for all of the various detection schemes discussed herein before, for several threshold values (both fixed and adaptive). The best results were found to be those for the AVM+I Re-Encode and the YI+MPMD detection schemes using adaptive thresholds as seen in FIGS. **4-5**; wherein FIG. **4** is a graph illustrating simulated miss-probability plots associated with a plurality of detection schemes using adaptive threshold values; and FIG. **5** is a graph illustrating simulated false alarm probability plots associated with a plurality of detection schemes using adaptive threshold values; and FIG. **5** is a graph illustrating simulated false alarm probability plots associated with a plurality of detection schemes using adaptive threshold values. For the three cases shown, the best balance of results was achieved with the AVM+1 Re-Encode approach.

Comparing these results with other known detection schemes, the present inventors recommend most preferably utilizing the AVM+1 Re-Encode AVM+1 Re Encode detection scheme with adaptive threshold value  $t_{adaptive}$ =7. This method was found to not only yield comparable  $p_{miss}$  and better  $P_{fa}$  curves as shown in Table 1 below, but to also require less computational complexity and storage requirements than other known schemes as seen in Table 2, also shown below. The present inventors also noted that because one known detection scheme uses a fixed threshold value for the simulations, the respective  $P_{fa}$  curves actually increase beyond the 0 dB point, whereas the inventive results remain flat (due to the re-encoding).

#### TABLE 1

Performance of Detection Mechanisms								
Detection Scheme	$\mathbf{P}_{miss}$ (at 0 dB channel $\mathbf{E}_{b}/\mathbf{N}_{0}$ )	$\mathbf{P}_{F\!A}$ (at 0 dB channel $\mathbf{E}_b/\mathbf{N}_0$ )						
AVM + 1 Re-Encode (t adaptive = 7)	1.8%	1.6%						
YI + MPMD (t adaptive = 5)	.66%	6%						
YI + MPMD (t adaptive = 7)	4%	.8%						
Known scheme 1:	1%	2% *increasing						
YI + MPMD (t = 22)		6						
Known scheme 2: YI + MPMD + SER	4%	.6%						

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Complexity of Detection Schemes									
Detection scheme	Adds/subtracts	Multiplies	Comparisons	Logical &	Storage				
AVM + 1 Re-Encode with adaptive threshold	700 adds	8	40 (bit)	432	164 (160 bit, 4 real)				
YI + MPMD	8192 subtracts	None	8220 (real)	None	2048 (1024 bit, 1024 real)				
YI + MPMD with adaptive threshold	8508 (316 adds, 8192 subtracts)	8	8220 (real)	None	2052 (1024 bit, 1028 real)				
YI + MPMD + SER	9728 (1536 adds, 8192 subtracts)	None	8380 (8220 real, 160 bit)	1728	2048 (1024 bit, 1024 real)				
Absolute Viterbi Metric	None	None	None	None	None				
AVM + 1 Re-Encode	384 adds	None	40 (bit)	432	160 bit				
YI + MPMD	8192 subtracts	None	8220 (real)	None	2048 (1024 bit, 1024				
YI	8192 subtracts	None	8192 (real)	None	1024 bit (path labels				
MPMD	None	None	28 (real)	None	1024 real (PMDs)				
AVM + YI	8192 subtracts	None	8192 (real)	None	1024 bit (path labels)				
AVM + YI + 1 Re- Encode	8576 (384 adds, 8192 subtracts)	None	8232 (8192 real, 40 bit)	432	1184 bit				
SER	1536 adds	None	160 (bit)	1728	160 bit				
YI + MPMD + SER	9728 (1536 adds,	None	8380 (8220	1728	2208 (1184 bit,				
	8192 subtracts)		real, 160 bit)		1024 real)				
Estimator	316 adds	8	None	None	4 real				
YI + MPMD with	8508 (316 adds,	8	8220 (real)	None	2052 (1024 bit,				
adaptive threshold	8192 subtracts)				1028 real)				
AVM + 1 Re-Encode with adaptive threshold	700 adds	8	40 (bit)	432	164 (160 bit, 4 real)				

This invention has been described in considerable detail in order to provide those skilled in the HS-SCCH detection scheme art with the information need to apply the novel principles and to construct and use such specialized components as are required. In view of the foregoing descriptions, it should be apparent that the present invention represents a significant departure from the prior art in construction and operation. However, while particular embodiments of the present invention have been described herein in detail, it is to be understood that various alterations, modifications and substitutions can be made therein without departing in any way from the spirit and scope of the present invention, as defined in the claims which follow.

What is claimed is:

1. A shared control channel (SCCH) detection method  $_{60}$  comprising the steps of:

- descrambling a plurality of bits of a received signal on the SCCH with a sequence generated by convolutional encoding and puncturing a user identification signal to produce a plurality of descrambled bits;
- estimating an amplitude value in a transceiver associated with the descrambled bits on the SCCH; and

65

scaling a desired threshold value by the estimated amplitude value to generate an adaptive threshold value there from, such that the adaptive threshold value is substantially robust to amplitude variations associated with at least one SCCH.

2. The method according to claim 1 wherein the step of estimating an amplitude value associated with the descrambled bits on the SCCH comprises estimating a desired metric associated with the received signal on the SCCH.

**3**. The method according to claim **2** wherein the step of scaling a desired threshold value by the estimated amplitude value to generate an adaptive threshold value there from, such that the adaptive threshold value is substantially robust to amplitude variations associated with at least one SCCH comprises normalizing the descrambled bits by the estimated desired metric.

4. The method according to claim 3 further comprising the step of passing the normalized descrambled bits to a Viterbi decoder, such that the desired adaptive threshold is substantially robust to amplitude variations associated with at least one SCCH.

10

**5**. The method according to claim **2**, wherein the step of estimating a desired metric associated with the descrambled bits on the SCCH comprises the steps of:

- determining absolute values associated with at least one data signal;
- accumulating the absolute values; and
- determining the estimated desired metric in response to the accumulated absolute values.
- 6. The method according to claim 5 further comprising the steps of:
  - decoding the descrambled bits via a Viterbi decoder and generating a decoded signal there from; and
  - detecting the decoded signal to determine whether an adaptive threshold condition associated with the adaptive threshold value has been met, wherein the adaptive <sup>15</sup> threshold value is substantially robust to amplitude variations associated with at least one SCCH.

7. The method according to claim 1 wherein the step of scaling a desired threshold value by the estimated amplitude value to generate an adaptive threshold value there from, such that the adaptive threshold value is substantially robust to amplitude variations associated with at least one SCCH, comprises the steps of:

setting a threshold value on the path metric difference 25 between the two largest merging paths at each state in a Viterbi decoding trellis; and scaling each pat metric difference threshold value by the estimated amplitude value to determine the presence of a surviving path.

**8**. A shared control channel (SCCH) detection method 5 comprising the steps of:

- descrambling a plurality of bits of a received signal on the SCCH with a sequence generated by convolutional encoding and puncturing a user identification signal to produce a plurality of descrambled bits;
- estimating an amplitude value in a transceiver associated with the descrambled bits; and
- scaling a desired threshold value by the estimated amplitude value to generate an adaptive threshold value for the SCCH.
- **9**. A shared control channel (SCCH) detection method comprising the steps of:
  - receiving a shared control channel vector in a transceiver, the shared control channel vector scrambled by a sequence generated by convolutional encoding and puncturing a user identification signal:
  - descrambled the shared control channel vector to produce a plurality of descrambled bits;

normalizing the descrambled bits; and

applying the normalized descrambled bits to a Viterbi decoder.

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